

Model Checking - II

Virendra Singh

Associate Professor

Computer Architecture and Dependable Systems Lab.

Dept. of Electrical Engineering
Indian Institute of Technology
Bombay

<http://www.ee.iitb.ac.in/~viren>

viren@ee.iitb.ac.in



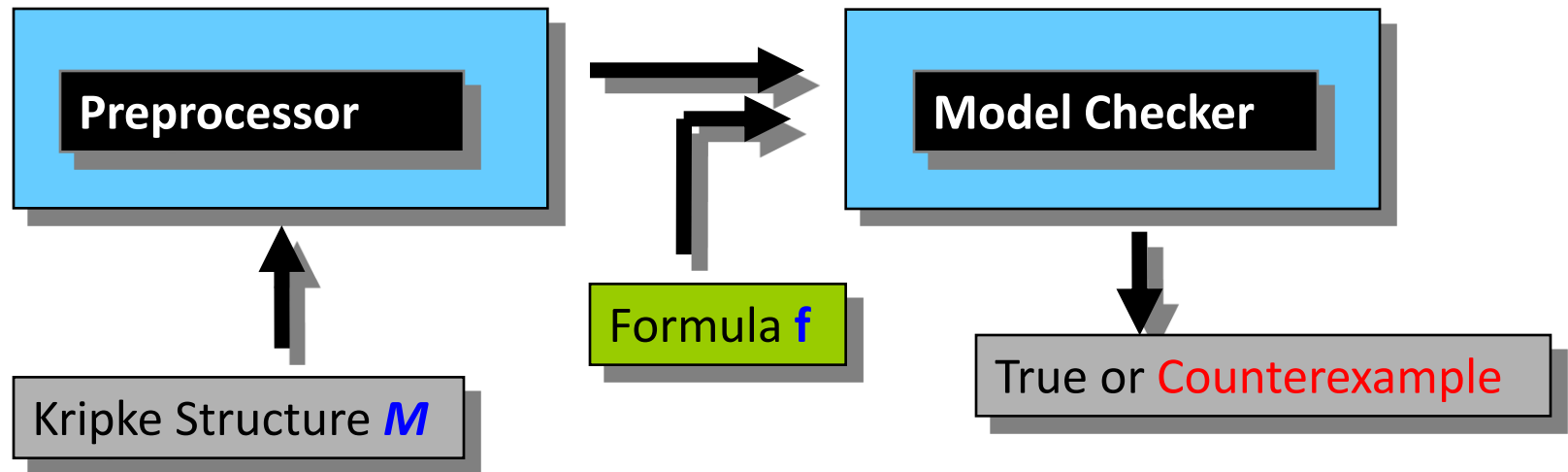
EE 709: Testing & Verification of VLSI Circuits

Lecture – 25 (Mar 01, 2012)

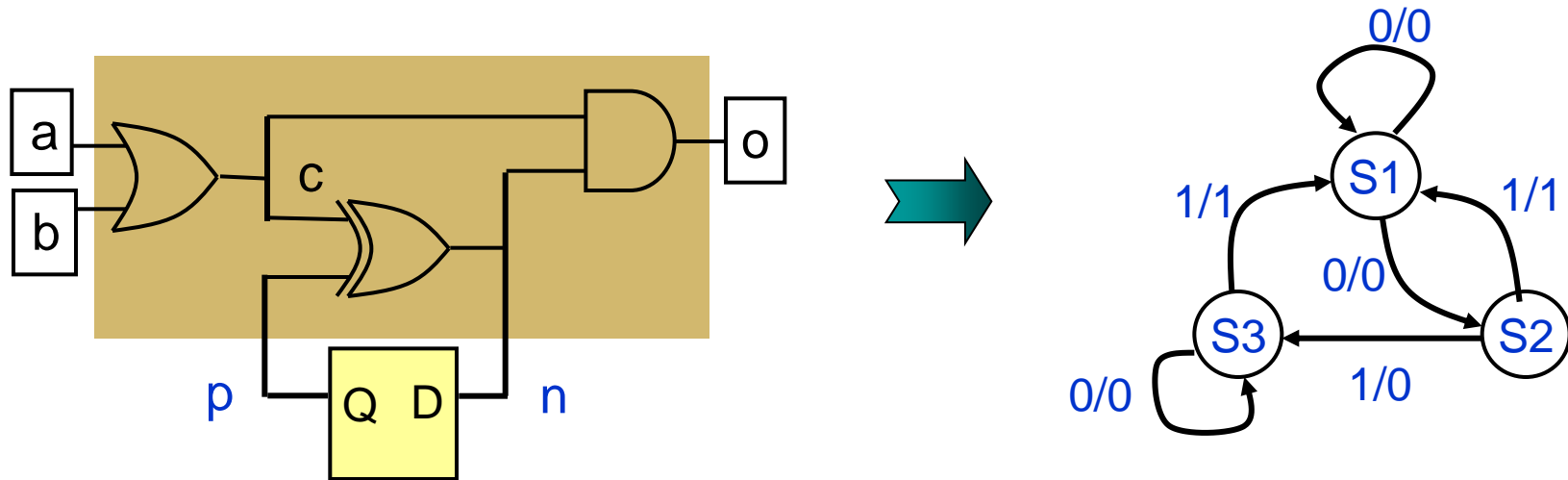
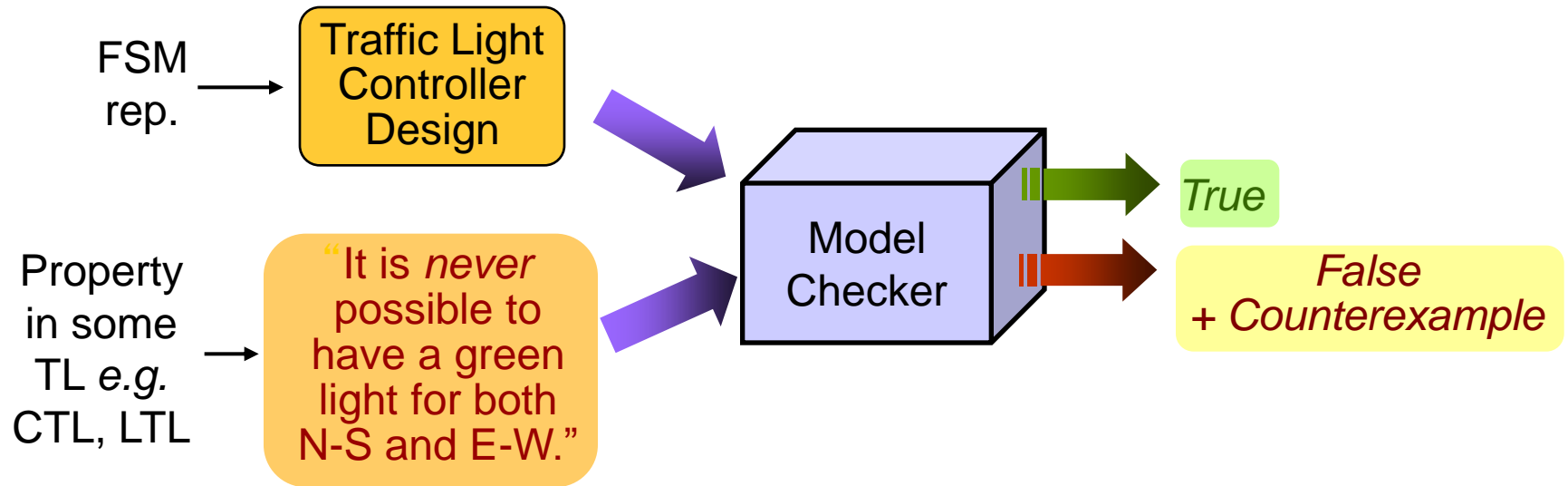
The Model Checking Problem

The Model Checking Problem (CE81):

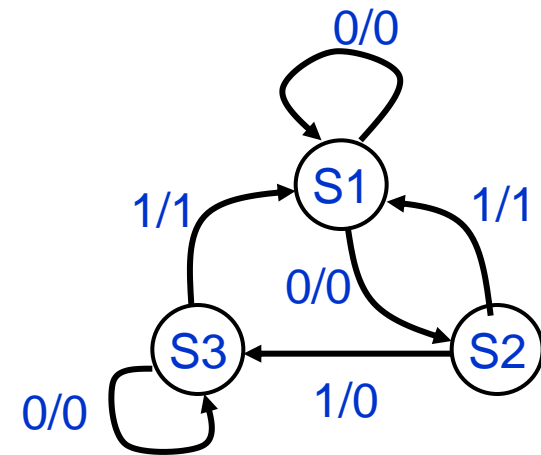
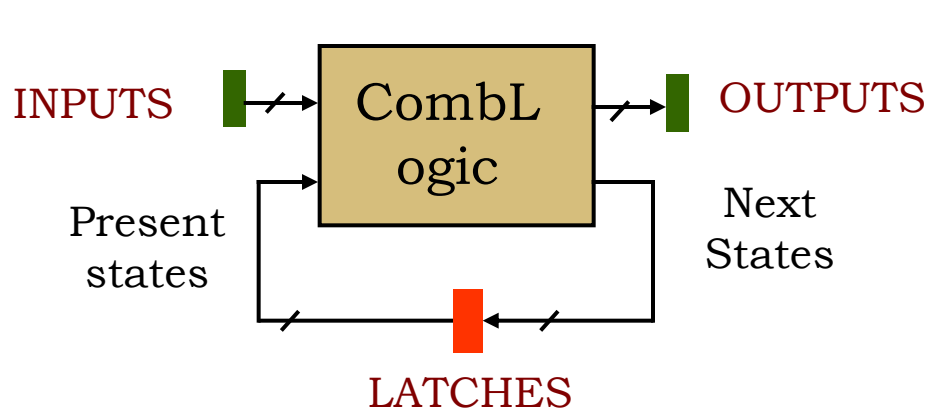
- Let M be a **Kripke structure** (i.e., state-transition graph).
- Let f be a **formula of temporal logic** (i.e., the specification).
- Find all states s of M such that $M, s \models f$



Temporal logic model checking



Finite State Machine (FSM)



State Transition Graph

Mealy FSM: $\langle I, S, \delta, S_0, O, \lambda \rangle$

- I : input alphabet
- S : finite, non-empty set of states
- $\delta : S \times I \rightarrow S$, next-state function
- $S^0 \subseteq S$: set of initial (reset) states
- O : output alphabet
- $\lambda : S \times I \rightarrow O$, output function

	$x = 0$	$x = 1$
$S1$	$S1,0$	$S2,1$
$S2$	$S1,0$	$S2,0$
$S3$	$S3,0$	$S1,1$

State Transition Table

3 Step Process

❖ Formal Specification

- Precise statement and property
- Environment constraint
- Logic: Temporal logic
- Automata, Labeled transition system

❖ Models

- Flexible model generation to specify design
- Fairness
- Transition system

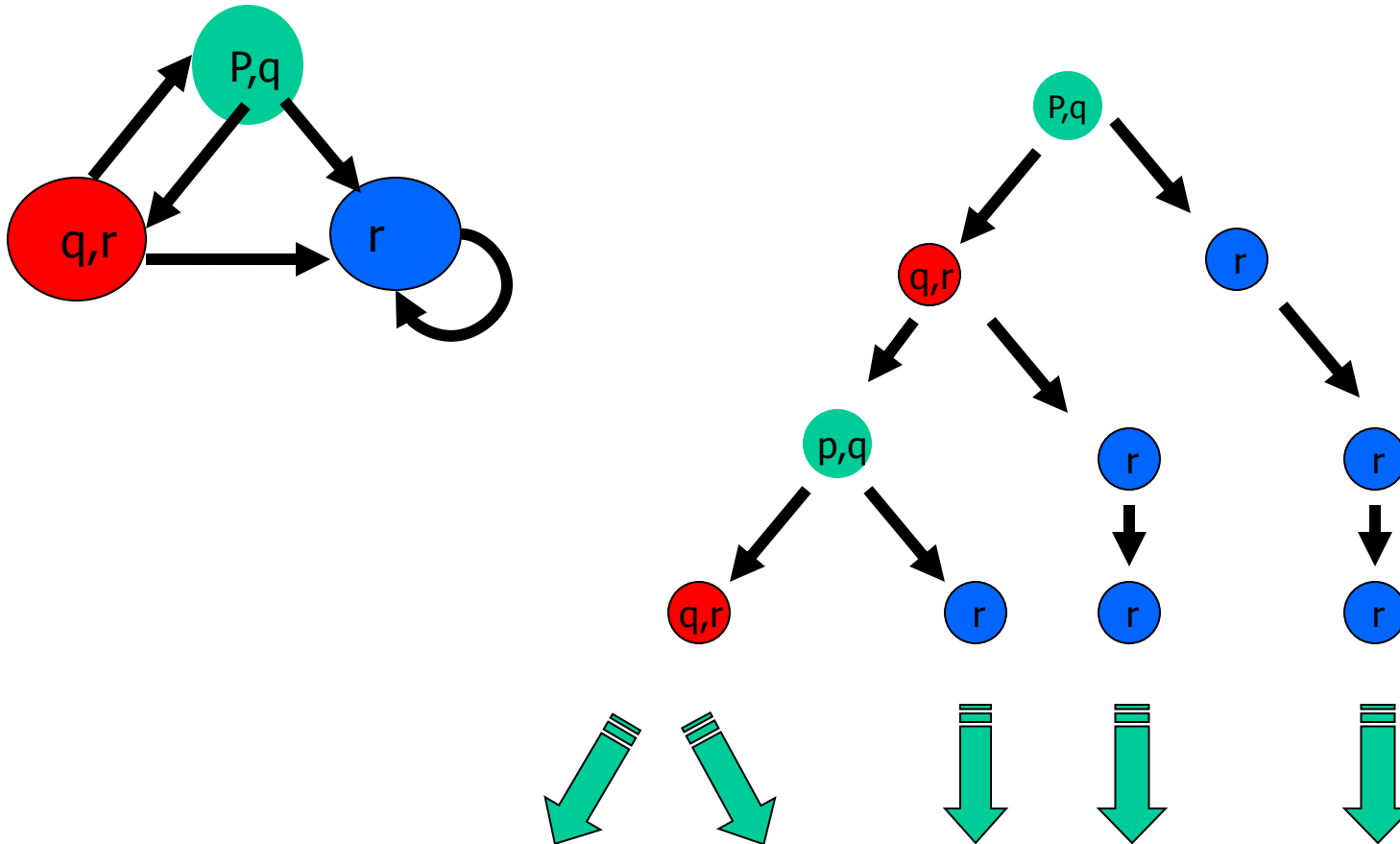
❖ Formal Verification

- Checking that model satisfy the property

Semantic of Finite State System

- ❖ Semantic associated with behaviour
- ❖ Branching Time Semantics
 - The tree of states obtained by unwinding the state machine transition graph
 - Possible choices are explicitly represented
- ❖ Linear Time Semantics
 - The set of all possible runs of the system
 - The set of infinite paths in SM

Computation Tree Logics



Formal Specification

- ❖ Describe unambiguously and precisely the expected behaviour of the design
- ❖ In general, a list of properties
- ❖ Includes, environmental constraints

Classification of Properties

- **Safety Property**

- (un) desirable things always (never) happen

- A bus arbiter never grants the requests to two masters
- Message received is message sent

- **Liveness (Progress) Property**

- desirable state eventually reached

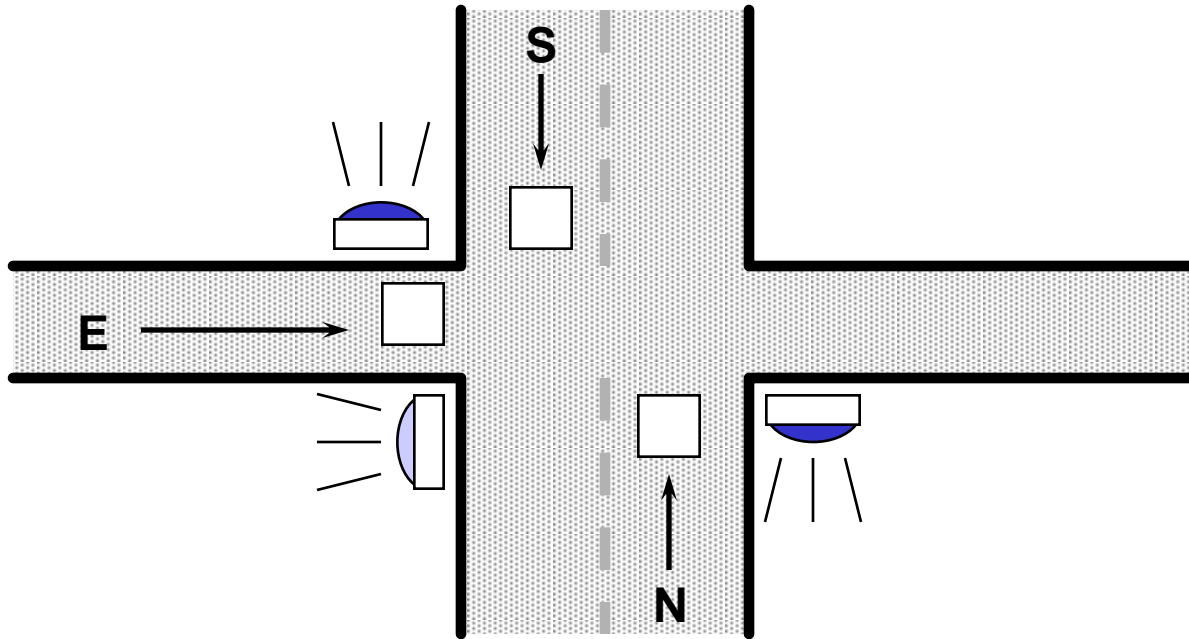
- Every bus request is eventually granted
- A car at a traffic light is eventually allowed to pass

- **Fairness Property**

- Desirable state repeatedly reached

- A request state and a grant state for each client must be visited infinitely often

Example: traffic light controller

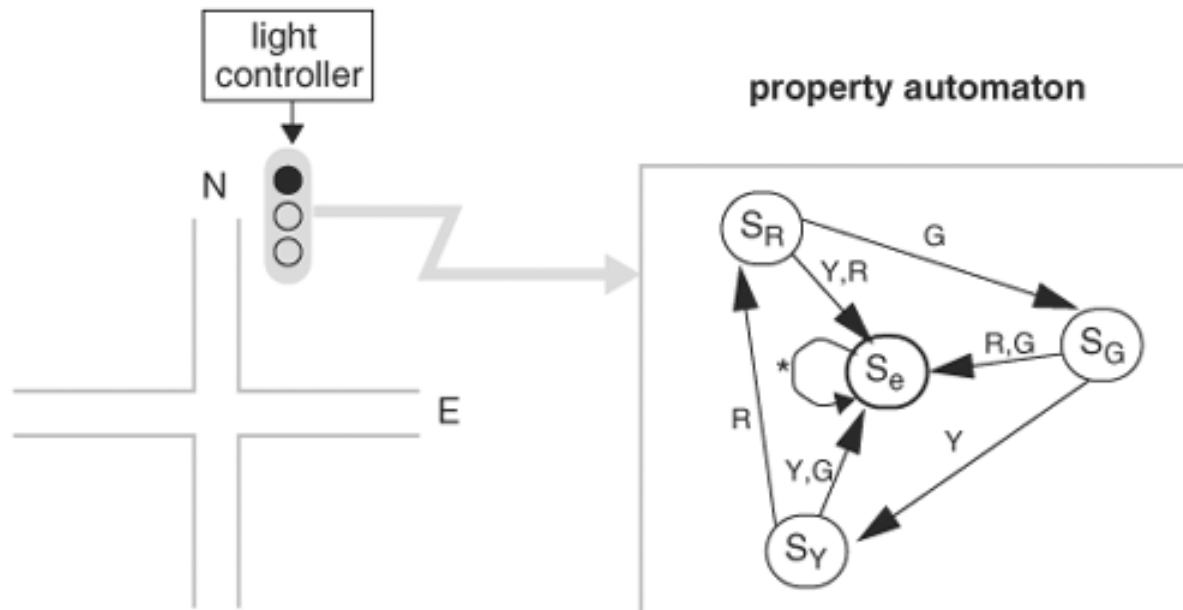


- Guarantee no collisions
- Guarantee eventual service

Property Specification

Properties for traffic light controller

- $P1 = (s1 \oplus w1) + (s2 \oplus w2)$
- Sequence R, G, Y, R, G, Y,



Thank you

