

# Bitcoin Transactions

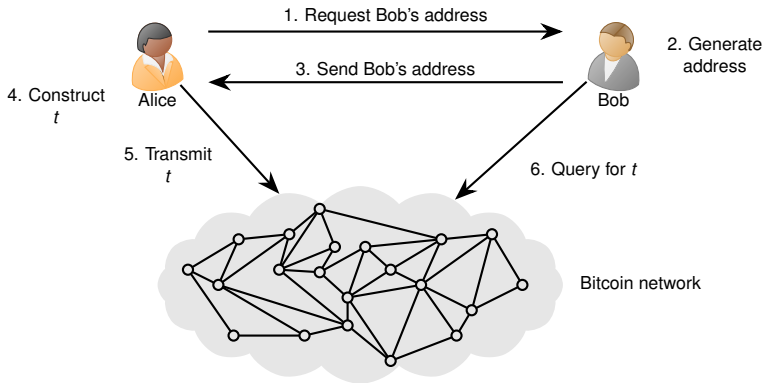
Saravanan Vijayakumaran  
sarva@ee.iitb.ac.in

Department of Electrical Engineering  
Indian Institute of Technology Bombay

January 11, 2024

# Bitcoin Transactions

# Bitcoin Payment Workflow



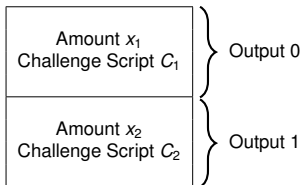
- Merchant Bob shares address out of band (not using Bitcoin P2P)
- Customer Alice broadcasts transaction  $t$  which pays the address
- Miners collect broadcasted transactions into a candidate block
- One of the candidate blocks containing  $t$  is mined
- Merchant waits for confirmations on  $t$  before providing goods

# Coinbase Transaction Format

## Block Format

Block Header
Number of Transactions $n$
Coinbase Transaction
Regular Transaction 1
Regular Transaction 2
⋮
Regular Transaction $n - 1$

## Coinbase Transaction

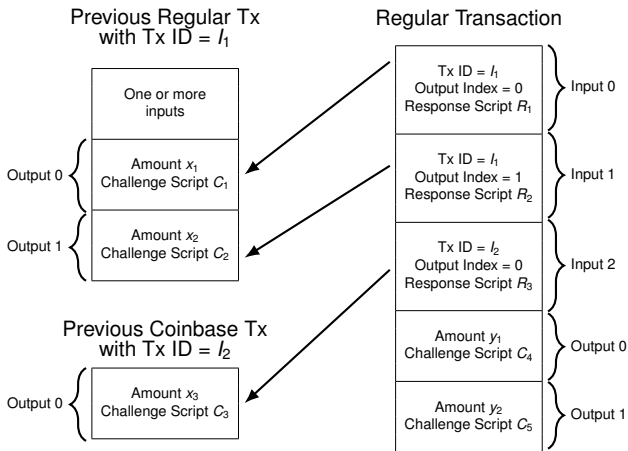


## Output Format

nValue
scriptPubkeyLen
scriptPubkey

- nValue contains number of satoshis locked in output
  - 1 Bitcoin =  $10^8$  satoshis
- scriptPubkey contains the challenge script
- scriptPubkeyLen contains byte length of challenge script

# Regular Transaction Format



## Input Format

hash
n
scriptSigLen
scriptSig
nSequence

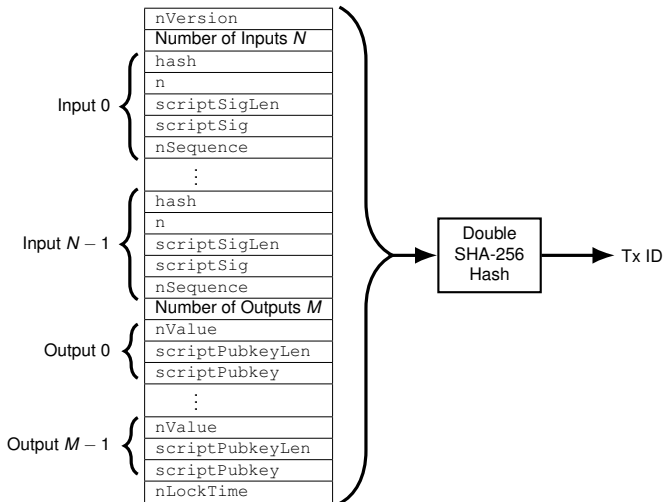
## Output Format

nValue
scriptPubkeyLen
scriptPubkey

- hash and n identify output being unlocked
- scriptSig contains the response script

# Transaction ID

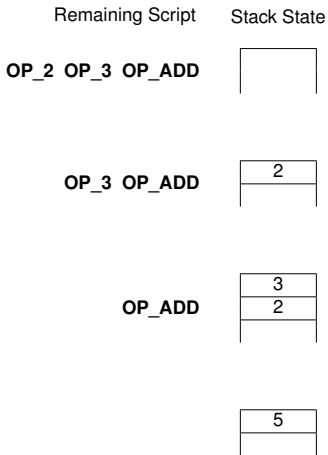
## Regular Transaction



# Bitcoin Scripting Language

# Script

- Forth-like stack-based language
- One-byte opcodes





# Challenge/Response Script Execution

Remaining Script

Stack State

**<Response Script>** **<Challenge Script>**



**<Challenge Script>**

$x_1$
$x_2$
$\vdots$
$x_n$

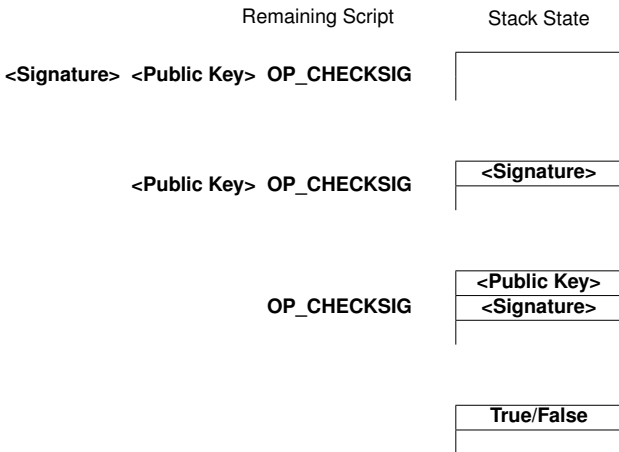
$y_1$
$y_2$
$\vdots$
$y_m$

Response is valid if top element  $y_1$  evaluates to True

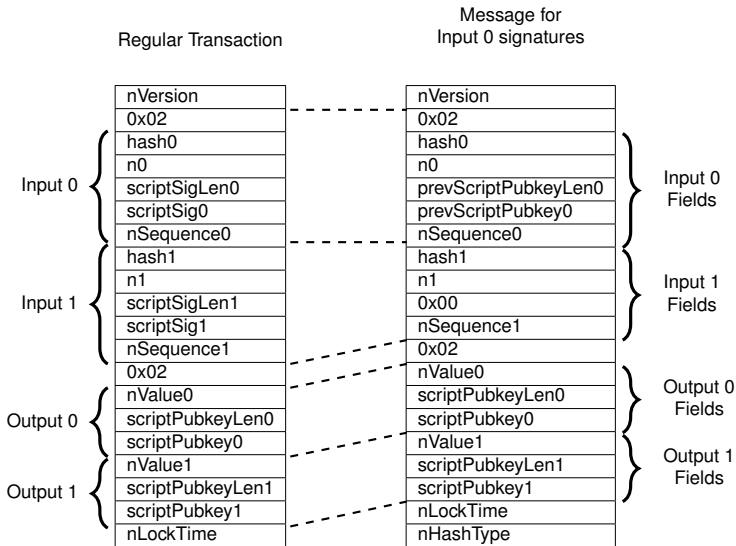


# Pay to Public Key

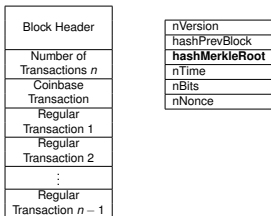
- Challenge script: **0x21 <Public Key> OP\_CHECKSIG**
- Response script: **<Signature>**



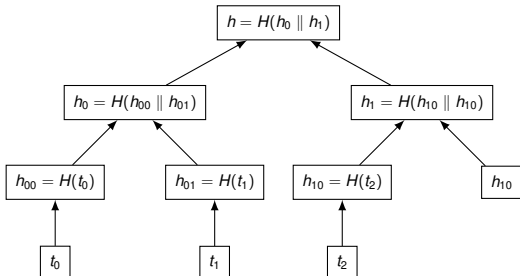
# Signatures Protect Transactions



# Transaction Merkle Root



- hashMerkleRoot contains root hash of transaction Merkle tree
- Modifying any transaction will modify the block header



## Key Takeaways

- Coinbase transactions have no inputs; outputs have challenge scripts
- Regular transaction inputs unlock previous outputs; outputs again have challenge scripts
- Scripts are expressed in a stack-based language
- Signatures prevent tampering of unconfirmed transactions

# Bitcoin Addresses

# Bitcoin Addresses

- To receive bitcoins, a challenge script needs to be specified
- Bitcoin addresses encode challenge scripts
- Example: 1EHNa6Q4Jz2uvNExL497mE43ikXhwF6kZm



- Bitcoin payment workflow (recap)
  - Merchant shares address out of band (not using Bitcoin P2P network)
  - Customer transmits transaction which pays the address
  - Merchant waits for transaction confirmations before providing goods/service



# Base58 Encoding

1EHNa6Q4Jz2uvNEXL497mE43ikXhwF6kZm



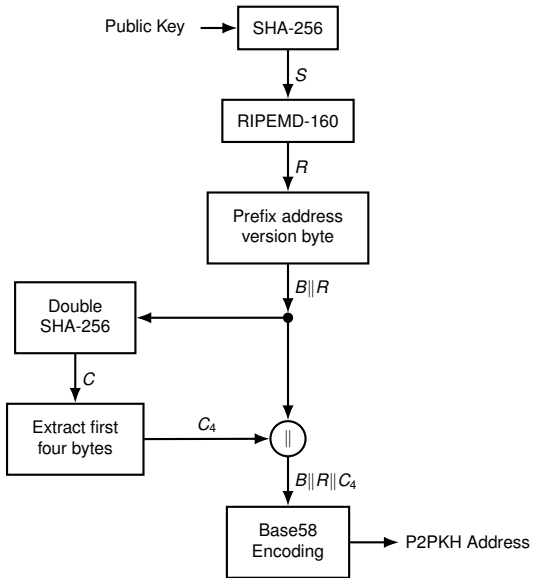
0091B24BF9F5288532960AC687ABB035127B1D28A50074FFE0

- Alphanumeric representation of bytestrings
- From 62 alphanumeric characters 0, O, I, l are excluded

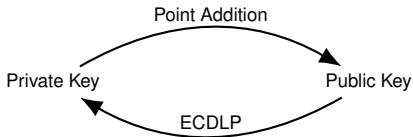
Ch	Int	Ch	Int	Ch	Int	Ch	Int	Ch	Int	Ch	Int	Ch	Int
1	0	A	9	K	18	U	27	d	36	n	45	w	54
2	1	B	10	L	19	V	28	e	37	o	46	x	55
3	2	C	11	M	20	W	29	f	38	p	47	y	56
4	3	D	12	N	21	X	30	g	39	q	48	z	57
5	4	E	13	P	22	Y	31	h	40	r	49		
6	5	F	14	Q	23	Z	32	i	41	s	50		
7	6	G	15	R	24	a	33	j	42	t	51		
8	7	H	16	S	25	b	34	k	43	u	52		
9	8	J	17	T	26	c	35	m	44	v	53		

- Given a bytestring  $b_n b_{n-1} \cdots b_0$ 
  - Encode each leading zero byte as a 1
  - Get integer  $N = \sum_{i=0}^{n-m} b_i 256^i$
  - Get  $a_k a_{k-1} \cdots a_0$  where  $N = \sum_{i=0}^k a_i 58^i$
  - Map each integer  $a_i$  to a Base58 character

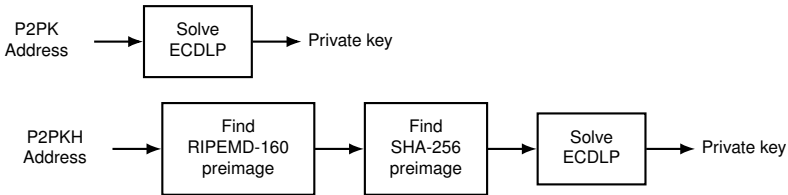
# Pay to Public Key Hash Address



# Why Hash the Public Key?

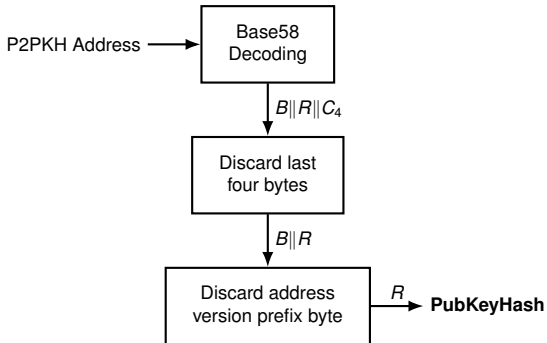


- ECDLP = Elliptic Curve Discrete Logarithm Problem
- ECDLP currently hard but no future guarantees
- Hashing the public key gives extra protection



# P2PKH Transaction

- Challenge script  
**OP\_DUP OP\_HASH160 <PubKeyHash> OP\_EQUALVERIFY  
OP\_CHECKSIG**



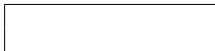
- Response script: **<Signature> <Public Key>**

# P2PKH Script Execution (1/2)

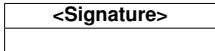
Remaining Script

Stack State

<Signature> <Public Key> OP\_DUP OP\_HASH160  
<PubKeyHash> OP\_EQUALVERIFY OP\_CHECKSIG



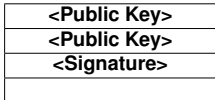
<Public Key> OP\_DUP OP\_HASH160  
<PubKeyHash> OP\_EQUALVERIFY OP\_CHECKSIG



OP\_DUP OP\_HASH160  
<PubKeyHash> OP\_EQUALVERIFY OP\_CHECKSIG



OP\_HASH160  
<PubKeyHash> OP\_EQUALVERIFY OP\_CHECKSIG



# P2PKH Script Execution (2/2)

Remaining Script

Stack State

**<PubKeyHash> OP\_EQUALVERIFY OP\_CHECKSIG**

<b>&lt;PubKeyHashCalc&gt;</b>
<b>&lt;Public Key&gt;</b>
<b>&lt;Signature&gt;</b>

**OP\_EQUALVERIFY OP\_CHECKSIG**

<b>&lt;PubKeyHash&gt;</b>
<b>&lt;PubKeyHashCalc&gt;</b>
<b>&lt;Public Key&gt;</b>
<b>&lt;Signature&gt;</b>

**OP\_CHECKSIG**

<b>&lt;Public Key&gt;</b>
<b>&lt;Signature&gt;</b>

**True/False**

--

## *m*-of-*n* Multi-Signature Scripts

- *m*-of-*n* multisig challenge script specifies *n* public keys

**m** <Public Key 1> ... <Public Key n> **n** OP\_CHECKMULTISIG

- Response script provides signatures created using **any** *m* out of the *n* private keys

**OP\_0** <Signature 1> ... <Signature m>.

- Example: *m* = 2 and *n* = 3

- Challenge script

**OP\_2** <PubKey1> <PubKey2> <PubKey3> **OP\_3** OP\_CHECKMULTISIG

- Response script

**OP\_0** <Sig1> <Sig2>

## 2-of-3 Multisig Script Execution

Remaining Script

OP\_0 <Sig1> <Sig2> OP\_2 <PubKey1>  
<PubKey2> <PubKey3> OP\_3 OP\_CHECKMULTISIG

OP\_2 <PubKey1>  
<PubKey2> <PubKey3> OP\_3 OP\_CHECKMULTISIG

OP\_CHECKMULTISIG

Stack State

--

<Sig2>
<Sig1>
<Empty Array>

3
<PubKey3>
<PubKey2>
<PubKey1>
2
<Sig2>
<Sig1>
<Empty Array>

True/False





# P2SH Multisig Script Execution (1/2)

Remaining Script

Stack State

OP\_0 <Sig1>  
<OP\_1 <PubKey1> <PubKey2> OP\_2 OP\_CHECKMULTISIG>  
OP\_HASH160 <RedeemScriptHash> OP\_EQUAL

--

<OP\_1 <PubKey1> <PubKey2> OP\_2 OP\_CHECKMULTISIG>  
OP\_HASH160 <RedeemScriptHash> OP\_EQUAL

<Sig1>
<Empty Array>

OP\_HASH160 <RedeemScriptHash> OP\_EQUAL

OP_1 <PubKey1> <PubKey2> OP_2 OP_CHECKMULTISIG
<Sig1>
<Empty Array>

<RedeemScriptHash> OP\_EQUAL

<RedeemScriptHashCalc>
<Sig1>
<Empty Array>

OP\_EQUAL

<RedeemScriptHash>
<RedeemScriptHashCalc>
<Sig1>
<Empty Array>

# P2SH Multisig Script Execution (2/2)

Remaining Script

Stack State

**OP\_1 <PubKey1> <PubKey2> OP\_2 OP\_CHECKMULTISIG**

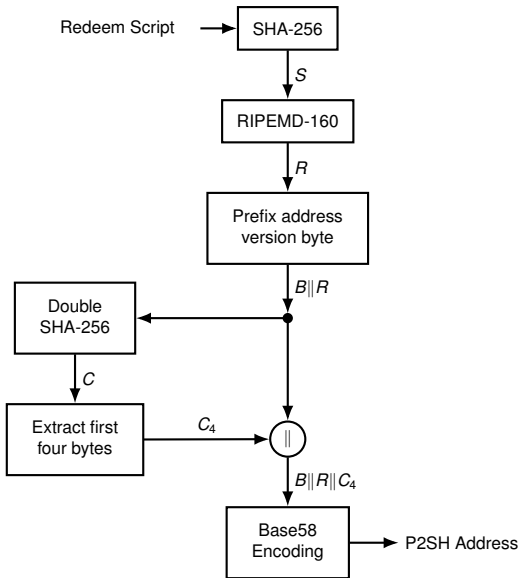
<b>&lt;Sig1&gt;</b>
<b>&lt;Empty Array&gt;</b>

**OP\_CHECKMULTISIG**

<b>2</b>
<b>&lt;PubKey2&gt;</b>
<b>&lt;PubKey1&gt;</b>
<b>1</b>
<b>&lt;Sig1&gt;</b>
<b>&lt;Empty Array&gt;</b>

<b>True/False</b>

# Pay to Script Hash Address



# Null Data Script

- Challenge script

**OP\_RETURN** <Data>

Length(<Data>)  $\leq$  80 bytes

- **OP\_RETURN** terminates script execution immediately
- No valid response script exists
  - Null data outputs are unspendable
  - Any bitcoins locked by a null data challenge script are lost forever
- Mainly used to timestamp data

# Pre-SegWit Standard Scripts

- Pay to Public Key (P2PK)
- Pay to Public Key Hash (P2PKH)
- *m-of-n* Multi-Signature (Multisig)
- Pay to Script Hash (P2SH)
- Null Data

## Key Takeaways

- Bitcoin addresses are shared over the Internet
- Transactions paying these addresses are broadcast on the Bitcoin network
- P2PKH addresses are obtained by hashing public keys
- Signatures created using private keys unlock P2PKH outputs
- P2SH addresses are obtained by hashing scripts
- Unlocking P2SH outputs requires both redeem script and valid response to it
- Null data scripts are for recording arbitrary data on the blockchain

## References

- Chapter 5 of *An Introduction to Bitcoin*, S. Vijayakumaran, [www.ee.iitb.ac.in/~sarva/bitcoin.html](http://www.ee.iitb.ac.in/~sarva/bitcoin.html)
- Bitcoin Script <https://en.bitcoin.it/wiki/Script>