## Matrix Computations EE636 Assignment 1

**Q** 1 Given an  $n \times n$  symmetric matrix A, calculate the number of flops required for finding:

- (a) LU factorization of A,
- (b) Cholesky decomposition.

 ${\bf Q}~{\bf 2}\,$  Calculate the flop count for forward substitution and backward substitution.

**Q** 3 Let  $A \in \mathbb{R}^{n \times n}$  be a (nonsingular) matrix whose leading principal submatrices are all nonsingular. Partition A as:

$$A = \begin{bmatrix} A_{11} & A_{12} \\ A_{21} & A_{22} \end{bmatrix}, \text{ where } A_{11} \in \mathbb{R}^{k \times k}, A_{12} \in \mathbb{R}^{k \times (n-k)}, A_{21} \in \mathbb{R}^{(n-k) \times k}, A_{22} \in \mathbb{R}^{(n-k) \times (n-k)}.$$

Show that there is exactly one matrix  $M \in \mathbb{R}^{(n-k) \times k}$  such that:

$$\begin{bmatrix} I_{\mathbf{k}} & 0\\ -M & I_{\mathbf{n}-\mathbf{k}} \end{bmatrix} \begin{bmatrix} A_{11} & A_{12}\\ A_{21} & A_{22} \end{bmatrix} = \begin{bmatrix} A_{11} & A_{12}\\ 0 & \tilde{A}_{22} \end{bmatrix}$$
(1)

Also calculate  $\tilde{A}_{22}$ .

 ${\bf Q}~{\bf 4}~{\rm Let}$ 

$$A = \begin{bmatrix} 2 & 2 & -4 \\ 1 & 1 & 5 \\ 1 & 3 & 6 \end{bmatrix} \text{ and } b = \begin{bmatrix} 10 \\ -2 \\ -5 \end{bmatrix}$$
(2)

Use Gaussian elimination with partial pivoting to find matrices L and U such that U is upper triangular, L is unit lower triangular with  $|l_{ij}| \leq 1$  for all i > j, and  $LU = \hat{A}$ , where  $\hat{A}$  can be obtained from A by making row interchanges. Use LU decomposition to solve the system Ax = b.

**Q** 5 Let A be positive definite. Then prove that A can be expressed in exactly one way as a product  $A = LDL^T$ , such that L is unit lower triangular, and D is a diagonal matrix whose main-diagonal entries are positive.

 $\mathbf{Q}$  6 For a symmetric matrix A, show that the following conditions are equivalent:

- (a)  $A \succeq 0$ .
- (b)  $A = U^T U$  for some matrix U.
- (c) All principal minors of A are nonnegative.

Note: A minor of of order k is principal if it is obtained by deleting (n - k) rows and the (n - k) columns with the same index numbers. For example: if one deletes the 3rd row and 4th row of a matrix  $A \in \mathbb{C}^{6\times 6}$ , then he/she should delete the 3rd and 4th column to get a principal minor of order 4.

**Q** 7 Consider the circuit in Figure 1.

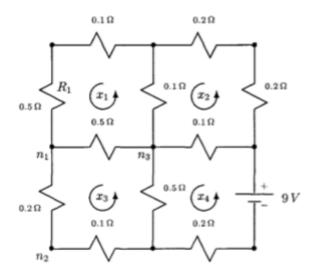


Figure 1: Electrical circuit with loop currents

- (a) Write down a linear system Ax = b of four equations for the four unknown loop currents.
- (b) Solve the system for x using the MATLAB command *linsolve*. Calculate the residual  $r = b A\hat{x}$ , where  $\hat{x}$  denotes your computed solution.
- (c) Using Ohm's law and the loop currents calculated in (b), find the voltage drops from the node labeled  $n_1$  to the nodes labeled  $n_2$  and  $n_3$ .
- **Q** 8 Suppose  $x(0:n) \in \mathbb{R}^{n+1}$ . A matrix  $V \in \mathbb{R}^{(n+1) \times (n+1)}$  of the form:

$$V = V(x_0, \dots, x_n) = \begin{bmatrix} 1 & 1 & \dots & 1 \\ x_0 & x_1 & \dots & x_n \\ \vdots & \vdots & \ddots & \vdots \\ x_0^n & x_1^n & \dots & x_n^n \end{bmatrix}$$
(3)

is said to be a Vandermonde matrix. Suppose  $x_i \neq x_j$  for  $i \neq j$ . This question is an illustration of how systems of the form  $V^T a = f = f(0:\mathbf{n})$  can be solved using  $O(\mathbf{n}^2)$  flops.

(a) Solving  $V^T a = f$  is equivalent to polynomial interpolation:

$$p(x) = \sum_{j=0}^{n} a_j x^j \tag{4}$$

where  $p(x_i) = f_i$  for i = 0: n. The first step in computing  $a_j$  of (5) is to calculate the Newton representation of the interpolating polynomial p:

$$p(x) = \sum_{k=0}^{n} c_k (\prod_{i=0}^{k-1} (x - x_i))$$
(5)

Devise an algorithm to determine the constants  $c_k$ .

(b) The next task is to generate the coefficients  $a_0, \ldots, a_n$  in (6) from the Newton representation coefficients  $c_0, \ldots, c_n$ . Define an iteration for the polynomials  $p_n(x), \ldots, p_0(x)$ which in turn provides a way to obtain coefficients  $a_i$  recursively. (c) Determine the flop count for the algorithm you devised combining (a) and (b).

**Q** 9 Consider the following system of equations:

$$\begin{bmatrix} 4 & 8 & -4 & -4 \\ 2 & 4 & 8 & -6 \\ -3 & -3 & 8 & -2 \\ -1 & 1 & 6 & -3 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \\ x_4 \end{bmatrix} = \begin{bmatrix} -8 \\ 10 \\ 7 \\ 7 \end{bmatrix}.$$

Use partial pivoting technique to solve for x.

**Q 10** Use complete pivoting technique to solve the following system of equations:

$$\begin{bmatrix} 3 & 4 & 3 \\ 1 & 5 & -1 \\ 6 & 3 & 7 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \end{bmatrix} = \begin{bmatrix} 10 \\ 7 \\ 15 \end{bmatrix}.$$

**Q 11** Consider the matrix  $F_{\mathbf{n}} \in \mathbb{C}^{\mathbf{n} \times \mathbf{n}}$  with the entries  $f_{\mathbf{jk}} = \omega_{\mathbf{n}}^{\mathbf{jk}}$ , where  $\mathbf{j}, \mathbf{k} \in \{0, 1, \dots, \mathbf{n-1}\}$  and  $\omega_{\mathbf{n}} = e^{-\frac{2\pi i}{\mathbf{n}}}$ .

(a) Suppose n = 2m, where m is a nutural number. Show that there exists a matrix P such that

$$F_{\mathbf{n}}P = \begin{bmatrix} F_{\mathbf{m}} & \Omega_{\mathbf{m}}F_{\mathbf{m}} \\ F_{\mathbf{m}} & -\Omega_{\mathbf{m}}F_{\mathbf{m}} \end{bmatrix},$$
  
where  $\Omega_{\mathbf{m}} = \operatorname{diag}(1, \omega_{\mathbf{n}}, \omega_{\mathbf{n}}^{2}, \dots, \omega_{\mathbf{n}}^{\mathbf{m}-1}).$ 

(b) Use the result found in part (a) to construct an algorithm to solve the system of equations

$$\bar{F}_{n}x = b$$

where  $\mathbf{n} = 2^{\mathbf{q}}$  for some  $\mathbf{q} \in \{0, 1, 2, ...\}$ .  $\overline{F}_{\mathbf{n}}$  is found by replacing each entry of  $F_{\mathbf{n}}$  by its complex conjugate.

(c) What will be the flop count for such an algorithm? Justify your answer.

**Q 12** For any  $v \in \mathbb{R}^n$  define the vectors  $v_+ = (v + E_n v)/2$  and  $v_- = (v - E_n v)/2$ . Suppose  $A \in \mathbb{R}^{n \times n}$  is symmetric and persymmetric. Show that if Ax = b then  $Ax_+ = b_+$  and  $Ax_- = b_-$ .

Note: 
$$E_{\mathbf{n}} := \begin{bmatrix} 0 & 0 & \dots & 0 & 1 \\ 0 & 0 & \dots & 1 & 0 \\ \vdots & \vdots & \ddots & \vdots & \vdots \\ 0 & 1 & \dots & 0 & 0 \\ 1 & 0 & \dots & 0 & 0 \end{bmatrix} \in \mathbb{R}^{\mathbf{n} \times \mathbf{n}}.$$