Bitcoin Transactions

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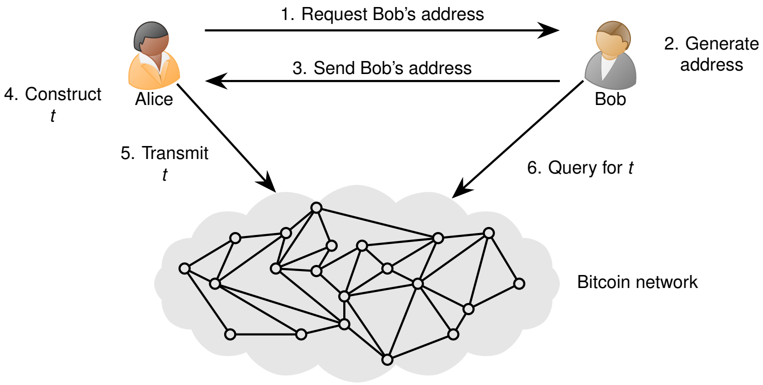
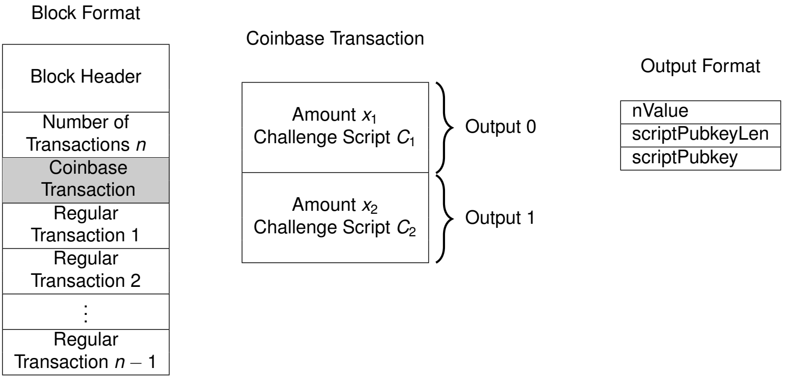
Bitcoin Payment Workflow



- Merchant Bob shares address out of band (not using Bitcoin P2P)
- Customer Alice broadcasts transaction *t* which pays the address
- Miners collect broadcasted transactions into a candidate block
- One of the candidate blocks containing *t* is mined
- Merchant waits for confirmations on *t* before providing goods

Coinbase Transaction Format



Block Format

Block Header

Number of Transactions n

Coinbase Transaction

Regular Transaction 1

Regular Transaction 2

$\vdots$

Regular Transaction $n-1$

Coinbase Transaction

Amount x_1
Challenge Script C_1

Output 0

Amount x_2
Challenge Script C_2

Output 1

Output Format

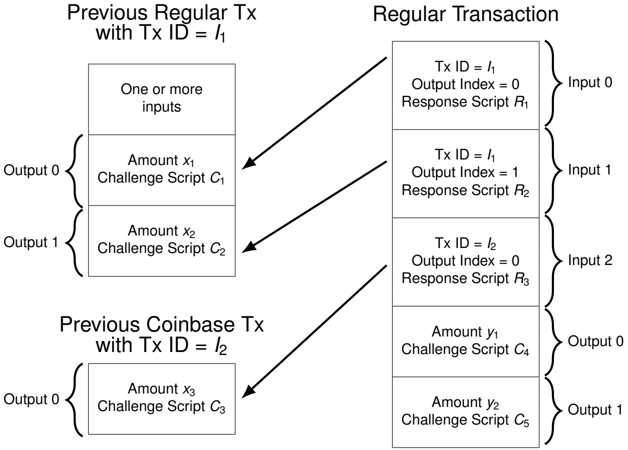
nValue

scriptPubkeyLen

scriptPubkey

- nValue contains number of satoshis locked in output
    - 1 Bitcoin = 10^8 satoshis
- scriptPubkey contains the challenge script
- scriptPubkeyLen contains byte length of challenge script

Regular Transaction Format



Previous Regular Tx with Tx ID = I_1

One or more inputs

Output 0: Amount x_1 Challenge Script C_1

Output 1: Amount x_2 Challenge Script C_2

Previous Coinbase Tx with Tx ID = I_2

Output 0: Amount x_3 Challenge Script C_3

Regular Transaction

Input 0: Tx ID = I_1 Output Index = 0 Response Script R_1

Input 1: Tx ID = I_1 Output Index = 1 Response Script R_2

Input 2: Tx ID = I_2 Output Index = 0 Response Script R_3

Output 0: Amount y_1 Challenge Script C_4

Output 1: Amount y_2 Challenge Script C_5

Input Format

hash
n
scriptSigLen
scriptSig
nSequence

Output Format

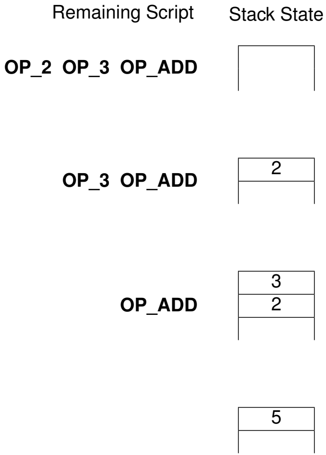
nValue
scriptPubkeyLen
scriptPubkey

- hash and n identify output being unlocked
- scriptSig contains the response script

Bitcoin Scripting Language

Script

- Forth-like stack-based language
- One-byte opcodes



Remaining Script	Stack State
OP_2 OP_3 OP_ADD	
OP_3 OP_ADD	2
OP_ADD	3 2
	5

Challenge/Response Script Execution

Remaining Script	Stack State
<Response Script> <Challenge Script>	
<Challenge Script>	x_1 x_2 ⋮ x_n
	y_1 y_2 ⋮ y_m

Response is valid if top element y_1 evaluates to True

Challenge Script Example

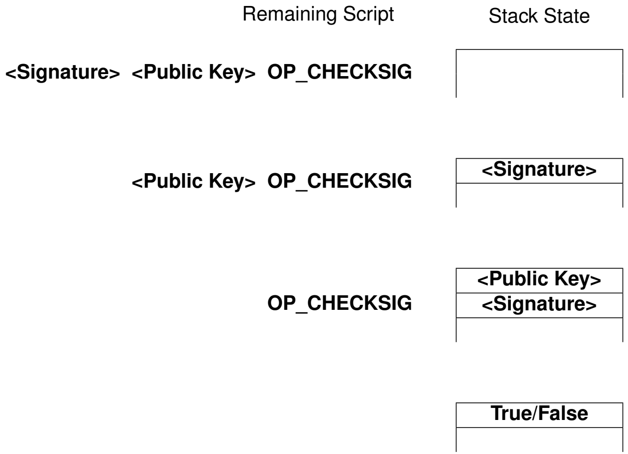
OP_HASH256 0x20 $\underbrace{\textbf{<256-bit string>}}_{\textbf{S}}$ **OP_EQUAL**

Remaining Script	Stack State
OP_HASH256 0x20 S OP_EQUAL	x
0x20 S OP_EQUAL	$H(x)$
OP_EQUAL	**S** $H(x)$
	0 or 1

Unsafe challenge script! Guess why?

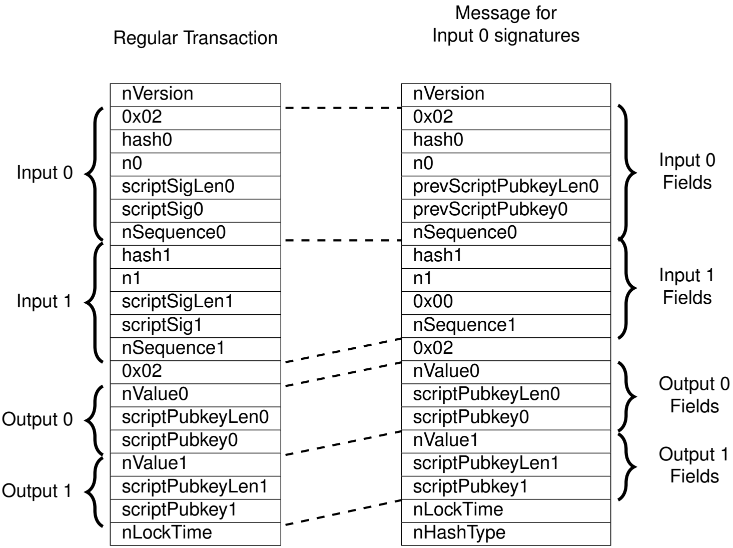
Pay to Public Key

- Challenge script: **0x21 <Public Key> OP_CHECKSIG**
- Response script: **<Signature>**



Remaining Script	Stack State
<Signature> <Public Key> OP_CHECKSIG	
<Public Key> OP_CHECKSIG	**<Signature>**
OP_CHECKSIG	**<Public Key>** **<Signature>**
	True/False

Signatures Protect Transactions



Regular Transaction

Input 0
- nVersion
- 0x02
- hash0
- n0
- scriptSigLen0
- scriptSig0
- nSequence0

Input 1
- hash1
- n1
- scriptSigLen1
- scriptSig1
- nSequence1
- 0x02

Output 0
- nValue0
- scriptPubkeyLen0
- scriptPubkey0

Output 1
- nValue1
- scriptPubkeyLen1
- scriptPubkey1
- nLockTime

Message for Input 0 signatures

- nVersion

Input 0 Fields
- 0x02
- hash0
- n0
- prevScriptPubkeyLen0
- prevScriptPubkey0
- nSequence0

Input 1 Fields
- hash1
- n1
- 0x00
- nSequence1
- 0x02

Output 0 Fields
- nValue0
- scriptPubkeyLen0
- scriptPubkey0

Output 1 Fields
- nValue1
- scriptPubkeyLen1
- scriptPubkey1
- nLockTime
- nHashType

Key Takeaways

- Coinbase transactions have no inputs; outputs have challenge scripts
- Regular transaction inputs unlock previous outputs; outputs again have challenge scripts
- Scripts are expressed in a stack-based language
- Signatures prevent tampering of unconfirmed transactions

Bitcoin Addresses

Bitcoin Addresses

- To receive bitcoins, a challenge script needs to be specified
- Bitcoin addresses encode challenge scripts
- Example: 1EHNa6Q4Jz2uvNExL497mE43ikXhwF6kZm

- Bitcoin payment workflow (recap)
  - Merchant shares address out of band (not using Bitcoin P2P network)
  - Customer transmits transaction which pays the address
  - Merchant waits for transaction confirmations before providing goods/service

Base58 Encoding

1EHNa6Q4Jz2uvNExL497mE43ikXhwF6kZm

$$\updownarrow$$

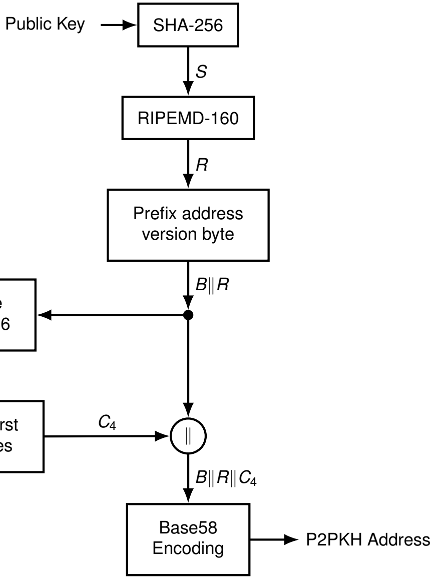
0091B24BF9F5288532960AC687ABB035127B1D28A50074FFE0

- Alphanumeric representation of bytestrings
- From 62 alphanumeric characters 0, O, I, l are excluded

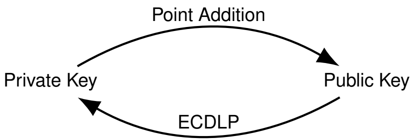
Ch	Int	Ch	Int	Ch	Int	Ch	Int	Ch	Int	Ch	Int	Ch	Int
1	0	A	9	K	18	U	27	d	36	n	45	w	54
2	1	B	10	L	19	V	28	e	37	o	46	x	55
3	2	C	11	M	20	W	29	f	38	p	47	y	56
4	3	D	12	N	21	X	30	g	39	q	48	z	57
5	4	E	13	P	22	Y	31	h	40	r	49		
6	5	F	14	Q	23	Z	32	i	41	s	50		
7	6	G	15	R	24	a	33	j	42	t	51		
8	7	H	16	S	25	b	34	k	43	u	52		
9	8	J	17	T	26	c	35	m	44	v	53		

- Given a bytestring $b_n b_{n-1} \cdots b_0$
  - Encode each leading zero byte as a 1
  - Get integer $N = \sum_{i=0}^{n-m} b_i 256^i$
  - Get $a_k a_{k-1} \cdots a_0$ where $N = \sum_{i=0}^{k} a_i 58^i$
  - Map each integer a_i to a Base58 character

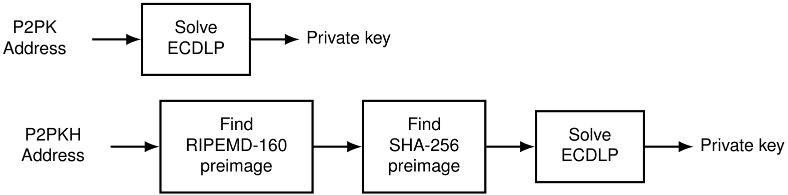
Pay to Public Key Hash Address



Why Hash the Public Key?

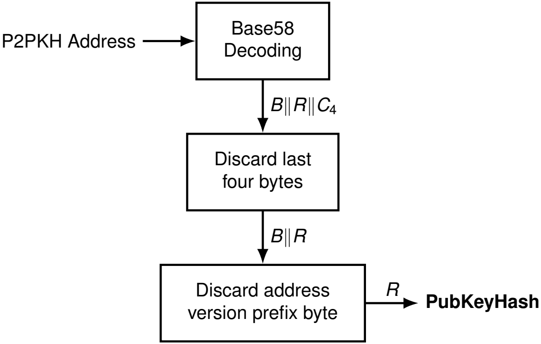


- ECDLP = Elliptic Curve Discrete Logarithm Problem
- ECDLP currently hard but no future guarantees
- Hashing the public key gives extra protection



P2PKH Transaction

- Challenge script
  OP_DUP OP_HASH160 <PubKeyHash> OP_EQUALVERIFY OP_CHECKSIG



- Response script: **<Signature> <Public Key>**

P2PKH Script Execution (1/2)

Remaining Script	Stack State
<Signature> <Public Key> OP_DUP OP_HASH160 <PubKeyHash> OP_EQUALVERIFY OP_CHECKSIG	
<Public Key> OP_DUP OP_HASH160 <PubKeyHash> OP_EQUALVERIFY OP_CHECKSIG	**<Signature>**
OP_DUP OP_HASH160 <PubKeyHash> OP_EQUALVERIFY OP_CHECKSIG	**<Public Key> <Signature>**
OP_HASH160 <PubKeyHash> OP_EQUALVERIFY OP_CHECKSIG	**<Public Key> <Public Key> <Signature>**

P2PKH Script Execution (2/2)

Remaining Script	Stack State
<PubKeyHash> OP_EQUALVERIFY OP_CHECKSIG	**<PubKeyHashCalc>** **<Public Key>** **<Signature>**
OP_EQUALVERIFY OP_CHECKSIG	**<PubKeyHash>** **<PubKeyHashCalc>** **<Public Key>** **<Signature>**
OP_CHECKSIG	**<Public Key>** **<Signature>**
	True/False

m-of-*n* Multi-Signature Scripts

- *m*-of-*n* multisig challenge script specifies *n* public keys

  m <Public Key 1> $\cdots$ <Public Key n> n OP_CHECKMULTISIG

- Response script provides signatures created using **any** *m* out of the *n* private keys

  OP_0 <Signature 1> $\cdots$ <Signature m>.

- Example: $m = 2$ and $n = 3$
  - Challenge script

    OP_2 <PubKey1> <PubKey2> <PubKey3> OP_3 OP_CHECKMULTISIG

  - Response script

    OP_0 <Sig1> <Sig2>

2-of-3 Multisig Script Execution

Remaining Script	Stack State
OP_0 <Sig1> <Sig2> OP_2 <PubKey1> <PubKey2> <PubKey3> OP_3 OP_CHECKMULTISIG	
OP_2 <PubKey1> <PubKey2> <PubKey3> OP_3 OP_CHECKMULTISIG	**<Sig2>** **<Sig1>** **<Empty Array>**
OP_CHECKMULTISIG	**3** **<PubKey3>** **<PubKey2>** **<PubKey1>** **2** **<Sig2>** **<Sig1>** **<Empty Array>**
	True/False

Pay to Script Hash Script

- Specify arbitrary scripts as payment destinations
- Challenge script

OP_HASH160 <RedeemScriptHash> OP_EQUAL

- Response script

<Response To Redeem Script> <Redeem Script>

- Example
  - 1-of-2 Multisig Challenge Script

    OP_1 <PubKey1> <PubKey2> OP_2 OP_CHECKMULTISIG

  - 1-of-2 Multisig Response Script

    OP_0 <Sig1> or **OP_0 <Sig2>**

  - P2SH Multisig challenge script

    OP_HASH160 <RedeemScriptHash> OP_EQUAL

  - P2SH Multisig response script

$$\underbrace{\textbf{OP\_0 <Sig1>}}_{\substack{\text{Response to}\\ \text{Redeem Script}}} \quad \underbrace{\textbf{OP\_1 <PubKey1> <PubKey2> OP\_2 OP\_CHECKMULTISIG}}_{\text{Redeem Script}}$$

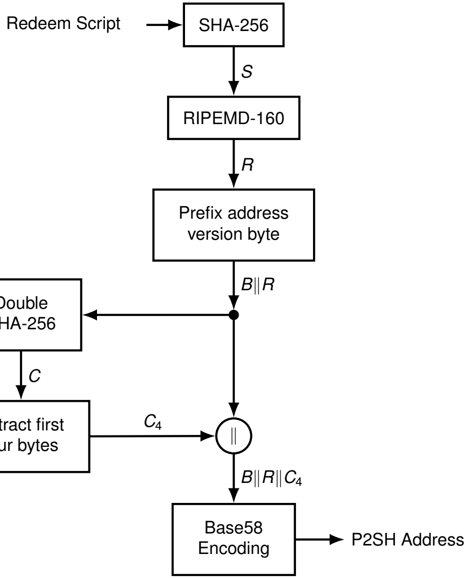
P2SH Multisig Script Execution (1/2)

Remaining Script	Stack State
OP_0 <Sig1> <OP_1 <PubKey1> <PubKey2> OP_2 OP_CHECKMULTISIG> OP_HASH160 <RedeemScriptHash> OP_EQUAL	
<OP_1 <PubKey1> <PubKey2> OP_2 OP_CHECKMULTISIG> OP_HASH160 <RedeemScriptHash> OP_EQUAL	<Sig1> <Empty Array>
OP_HASH160 <RedeemScriptHash> OP_EQUAL	OP_1 <PubKey1> <PubKey2> OP_2 OP_CHECKMULTISIG <Sig1> <Empty Array>
<RedeemScriptHash> OP_EQUAL	<RedeemScriptHashCalc> <Sig1> <Empty Array>
OP_EQUAL	<RedeemScriptHash> <RedeemScriptHashCalc> <Sig1> <Empty Array>

P2SH Multisig Script Execution (2/2)

Remaining Script	Stack State
OP_1 <PubKey1> <PubKey2> OP_2 OP_CHECKMULTISIG	**<Sig1>** **<Empty Array>**
OP_CHECKMULTISIG	**2** **<PubKey2>** **<PubKey1>** **1** **<Sig1>** **<Empty Array>**
	True/False

Pay to Script Hash Address



Null Data Script

- Challenge script

  `OP_RETURN <Data>`

  Length(**`<Data>`**) $\leq$ 80 bytes
- **`OP_RETURN`** terminates script execution immediately
- No valid response script exists
  - Null data outputs are unspendable
  - Any bitcoins locked by a null data challenge script are lost forever
- Mainly used to timestamp data

Pre-SegWit Standard Scripts

- Pay to Public Key (P2PK)
- Pay to Public Key Hash (P2PKH)
- *m*-of-*n* Multi-Signature (Multisig)
- Pay to Script Hash (P2SH)
- Null Data

Key Takeaways

- Bitcoin addresses are shared over the Internet
- Transactions paying these addresses are broadcast on the Bitcoin network
- P2PKH addresses are obtained by hashing public keys
- Signatures created using private keys unlock P2PKH outputs
- P2SH addresses are obtained by hashing scripts
- Unlocking P2SH outputs requires both redeem script and valid response to it
- Null data scripts are for recording arbitrary data on the blockchain

References

- Chapter 5 of *An Introduction to Bitcoin*, S. Vijayakumaran, `www.ee.iitb.ac.in/~sarva/bitcoin.html`