

1. [5 points] For a negligible function negl , prove that $p(n)\text{negl}(n)$ is also negligible for any positive polynomial p .
2. [5 points] Let $G : \{0, 1\}^n \rightarrow \{0, 1\}^{l(n)}$ be a pseudorandom generator with expansion factor $l(n) > n$. **Prove or disprove** that the following functions are pseudorandom generators where $s \in \{0, 1\}^n$ and s_i is the i th bit of s . The \parallel denotes the string concatenation operator.
 - (a) $G_1(s) = G(s)\parallel 0$.
 - (b) $G_2(s) = G(s_1, s_2, \dots, s_{|s|-1})\parallel s_{|s|}$.
 - (c) $G_3(s) = G(s)\parallel 0$.
3. [10 points] A private-key encryption scheme $\Pi = (\text{Gen}, \text{Enc}, \text{Dec})$ is **EAV-secure**, if for all PPT adversaries \mathcal{A} there is a negligible function negl such that, for all n ,

$$\Pr [\text{PrivK}_{\mathcal{A}, \Pi}^{\text{eav}}(n) = 1] \leq \frac{1}{2} + \text{negl}(n).$$

Let $\text{out}_{\mathcal{A}}(\text{PrivK}_{\mathcal{A}, \Pi}^{\text{eav}}(n, b))$ denote the output b' of \mathcal{A} when m_b is encrypted. Suppose that a private-key encryption scheme $\Pi = (\text{Gen}, \text{Enc}, \text{Dec})$ is EAV-secure.

Prove that for all PPT adversaries \mathcal{A} there is a negligible function negl such that, for all n ,

$$|\Pr [\text{out}_{\mathcal{A}}(\text{PrivK}_{\mathcal{A}, \Pi}^{\text{eav}}(n, 0)) = 1] - \Pr [\text{out}_{\mathcal{A}}(\text{PrivK}_{\mathcal{A}, \Pi}^{\text{eav}}(n, 1)) = 1]| \leq \text{negl}(n).$$