EE 720: Introduction to Number Theory and Cryptography (Autumn 2026)
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Assignment 1: 20 points — Date: August 6, 2026

1. [5 points] Consider the Vigenère cipher where the adversary knows that the key length is 100 characters. Let $S = \{0, 1, 2, \ldots, 25\}$. The key generation algorithm Gen generates the key $\mathbf{k} = k_0 k_1 k_2 \cdots k_{99}$ uniformly from the set S^{100}.

   Let $\mathcal{M} = \{0, 1, \ldots, 25\}^*$, i.e. the set of all finite length strings from the set $\{0, 1, \ldots, 25\}$. The encryption of a message $\mathbf{m} = m_0 m_1 \cdots m_{n-1} \in S^n$ is given by $\mathbf{c} = c_0 c_1 \cdots c_{n-1} \in S^n$ where $c_i = m_i + k_{i \bmod 100} \bmod 26$.

   Since the size of the key space is smaller than the size of the message space, this form of the Vigenère cipher is **not** perfectly secret. Then there must exist an adversary in the perfect indistinguishability experiment $\mathsf{PrivK}^{\mathrm{eav}}_{\mathcal{A},\Pi}$ that succeeds with a probability greater than $\frac{1}{2}$. Find such an adversary $\mathcal{A}$.

2. [10 points] When using the one-time pad with the key $k = 0^l$, we have $\mathtt{Enc}_k(m) = k \oplus m = m$ and the message is sent in the clear. It has therefore been suggested to modify the one-time pad by only encrypting with $k \neq 0^l$ (i.e., to have $\mathtt{Gen}$ choose k uniformly from the set of nonzero keys of length l).

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3. [5 points] For negligible functions negl_1 and negl_2, prove that $p_1(n)\mathsf{negl}_1(n) + p_2(n)\mathsf{negl}_2(n)$ is also negligible for any positive polynomials p_1, p_2.