

1. Suppose $F : \{0, 1\}^n \times \{0, 1\}^n \mapsto \{0, 1\}^n$ is a length-preserving pseudorandom function. **Using only F , give constructions** of schemes which satisfy the following properties. You can use results discussed in class without proof.
 - (a) [2 points] A CPA-secure encryption scheme for messages of length n which is not CCA-secure. Describe the algorithms (**Gen**, **Enc**, **Dec**).
 - (b) [2 points] A CPA-secure encryption scheme for messages of length $3n$. Describe the algorithms (**Gen**, **Enc**, **Dec**).
 - (c) [2 points] A secure MAC for messages of length $3n$. Describe the algorithms (**Gen**, **Mac**, **Vrfy**).
 - (d) [4 points] A CCA-secure encryption scheme for messages of length $3n$. Describe the algorithms (**Gen**, **Enc**, **Dec**).
2. [10 points] Alice has a length-preserving pseudorandom function $F : \{0, 1\}^n \times \{0, 1\}^n \rightarrow \{0, 1\}^n$. She wants to authenticate messages of length $2n$. Let $m \in \{0, 1\}^{2n}$ denote the message. Let $m_1 \in \{0, 1\}^n$ denote the first n bits of m and let $m_2 \in \{0, 1\}^n$ denote the last n bits of m . Alice uses the MAC scheme $\Pi = (\mathbf{Gen}, \mathbf{Mac}, \mathbf{Vrfy})$ where:
 - **Gen**: Key k is chosen uniformly from $\{0, 1\}^n$.
 - **Mac**: The message space $\mathcal{M} = \{0, 1\}^{2n}$. Given key k , the tag corresponding to $m = (m_1, m_2) \in \{0, 1\}^{2n}$ is given by a $2n$ -bit string t as follows.

$$\mathbf{Mac}_k(m) = t := \langle t_1, t_2 \rangle = \langle F_k(m_1), F_k(m_1 \oplus m_2) \rangle.$$

- **Vrfy**: Given key k , message $m = (m_1, m_2)$ and tag $t = \langle t_1, t_2 \rangle$, the verifier recomputes the tag on the message

$$t' := \langle t'_1, t'_2 \rangle = \langle F_k(m_1), F_k(m_1 \oplus m_2) \rangle.$$

If $t = t'$, then the verifier outputs 1 and 0 otherwise.

Prove that Alice's MAC scheme is **not** secure.